CS 331, Fall 2025

Lecture 18 (10/29)

Today: - PSD matrices
- Convexity in $\mathbb{R}^d$
- Linear regression
- Algos for LP

PSD matrices (Part VI, Section 4.2)

SVD $A = U \Sigma V^T$, $\Sigma = \text{diag}(\sigma)$ ← nonnegative

$$A = \sum_{i \in [d]} \sigma_i u_i v_i^T$$

u_i ← column of U; v_i ← column of V

Interpretation: each unit of $v_i \to \sigma_i$ units of u_i

$$Ax = U \Sigma V^T V c = U \Sigma c = \sum_{i \in [d]} \sigma_i c_i u_i$$

$Vc = \sum_{i \in [d]} c_i v_i$; V: orthonormal

Spectral theorem: $\forall$ symmetric $M \in \mathbb{R}^{d \times d}$

$$M = U \Lambda U^\top \qquad U^\top U = I$$

"orthonormal"

$$= \underbrace{\sum_{i \in [d]} \lambda_i u_i u_i^\top}_{\text{eigendecomposition}}$$

$$U = \begin{pmatrix} u_1 & u_2 & \cdots & u_d \end{pmatrix}$$

$$\Lambda = \begin{pmatrix} \lambda_1 & & & 0 \\ & \lambda_2 & & \\ & & \ddots & \\ 0 & & & \lambda_d \end{pmatrix}$$

Compare to SVD:

$$M = U \Sigma V^\top$$

$$= \sum_{i \in [d]} \sigma_i u_i v_i^\top$$

idea: if $\lambda_i < 0$

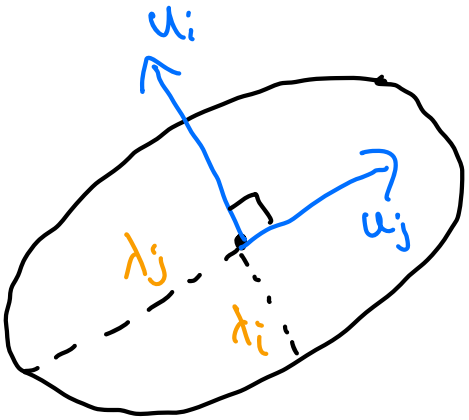
$\Rightarrow \sigma_i = -\lambda_i$

$v_i = -u_i$

If all $\lambda_i \geq 0$ already, SVD = eigendecomp.

We call such M PSD

"Every PSD matrix is an ellipse"



We say (u, λ) is eigvec/eigval of M if: $Mu = \lambda u$

"eigpair"

Claim: if $M = U \Lambda U^T = \sum_{i \in [d]} \lambda_i u_i u_i^T$

then (λ_j, u_j) is eigpair $\forall j \in [d]$

Proof: $\left(\sum_{i \in [d]} \lambda_i u_i u_i^T \right) u_j = \lambda_j u_j$

$\lambda_1 u_1 \underbrace{u_1^T u_j}_{=0} + \dots + \lambda_j u_j \underbrace{u_j^T u_j}_{=1} + \dots$

Main Claim: Let M be symmetric, $d \times d$.

It's PSD iff $v^T M v \geq 0 \quad \forall v \in \mathbb{R}^d$

Proof: Let $M = \sum_{i \in [d]} \lambda_i u_i u_i^T$

M not PSD $\Rightarrow \lambda_i < 0$ for some i

$\Rightarrow u_i^T M u_i = u_i^T (\lambda_i u_i) < 0$

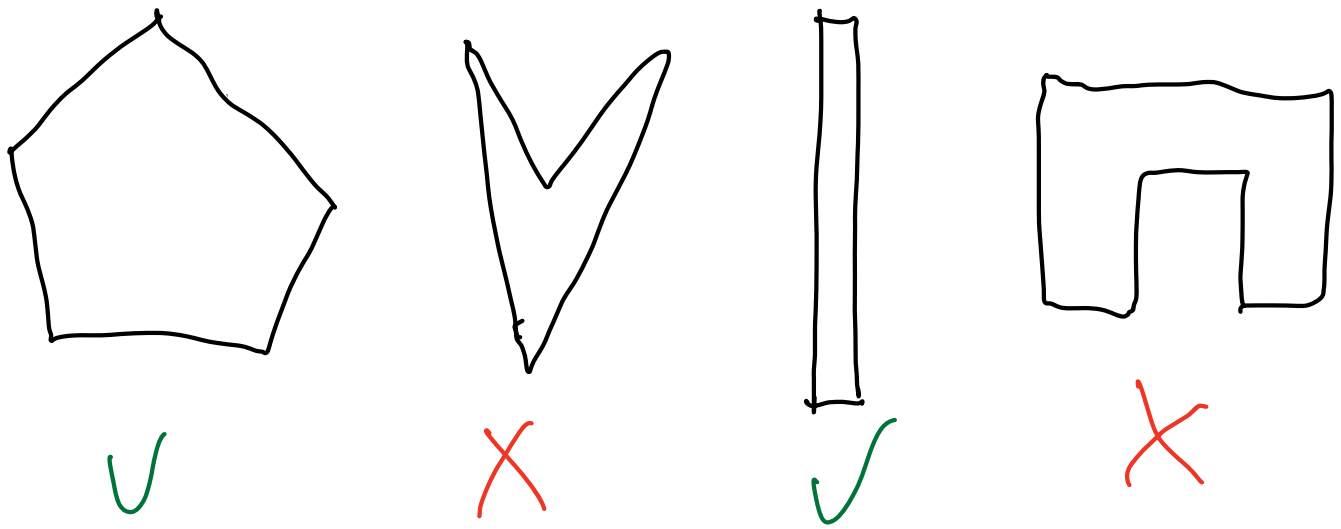
M is PSD $\Rightarrow v^T \left(\sum_{i \in [d]} \lambda_i u_i u_i^T \right) v$

$$= \sum_{i \in [d]} \underbrace{\lambda_i}_{\geq 0} \underbrace{\left(u_i^T v \right)^2}_{\geq 0} \geq 0.$$

e.g. M PSD means $M_{ii} \geq 0$: let $v = e_i$

Convexity in $\mathbb{R}^d$ (Part VI, Section 5.1)

Let $\chi \subseteq \mathbb{R}^d$. We say χ is <u>convex</u> if it contains all lines:



✓ ✗ ✓ ✗

We say $f : \chi \to \mathbb{R}$ is convex if: $\chi \subseteq \mathbb{R}^d$

- χ convex
- All <u>1-D restrictions</u> of f are convex (stay at/below the line)

Two examples so far:

Section 2 • $\chi = \mathbb{R}$ • $d = 1$

Section 5.3: Linear programming

Generalizes maxflow:

$$\max_{f \in \mathbb{R}^E} \sum_{v:(s,v)\in E} f_{(s,v)}$$

$$0 \le f_e \le c_e \quad \forall e \in E$$

$$\sum_{(v,u)\in E} f_{(v,u)} - \sum_{(u,v)\in E} f_{(u,v)} = 0 \quad \forall v \in V,\ v \notin \{s,t\}$$

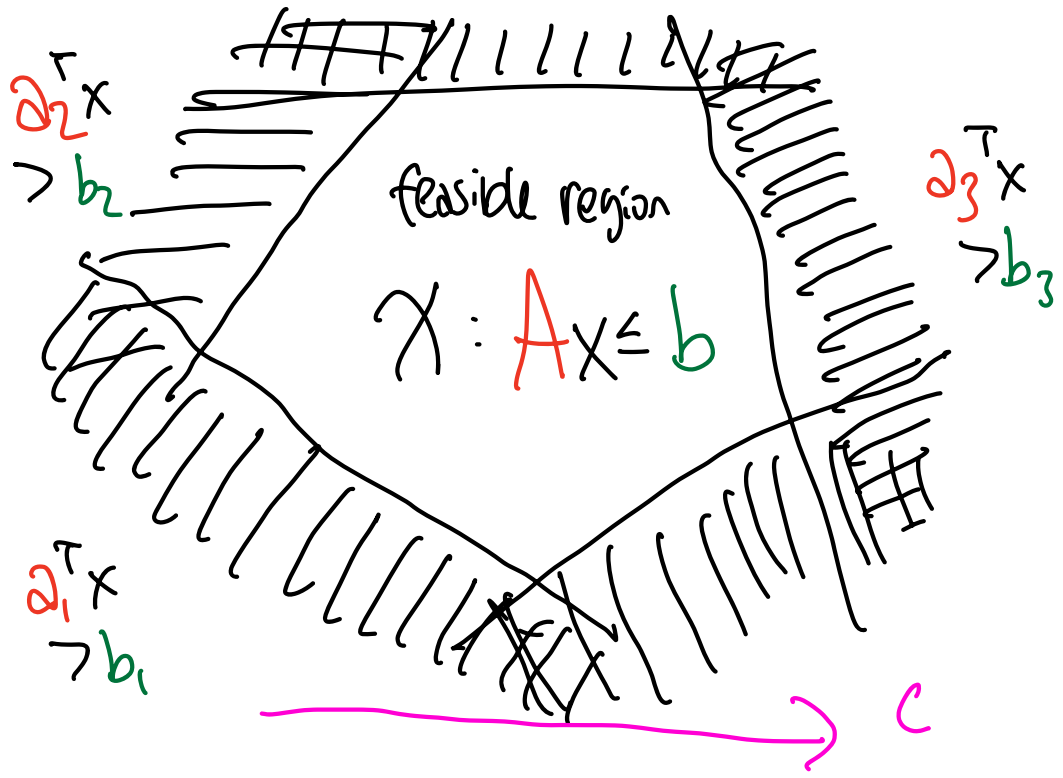
Linear program:

- $\chi = \{x \in \mathbb{R}^d \mid Ax \le b\}$

  * equality $a^T x = b$ is $a^T x \le b$, $-a^T x \le -b$

- $f(x) = c^T x = \sum_{i \in [d]} c_i x_i$

Pictorially:



f is convex:

$$f((1-\lambda)x + \lambda y) = c^T((1-\lambda)x + \lambda y)$$

$$= (1-\lambda)c^T x + \lambda c^T y$$

"at the line" $= (1-\lambda)f(x) + \lambda f(y)$ ✓

χ is convex: if $a_i^T x \leq b_i$

$a_i^T y \leq b_i$ $\forall i \in [n]$

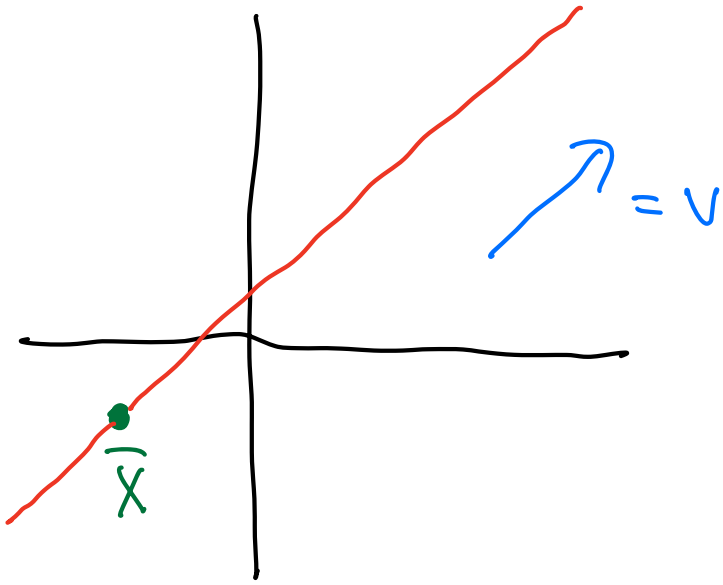
then $a_i^T((1-\lambda)x + \lambda y) \leq b_i$ $\forall i \in [n]$

Let $f: \mathcal{X} \to \mathbb{R}$ be convex, so:

$f_{\bar{x},v}(t)$

$= f(\bar{x} + tv)$

is convex.



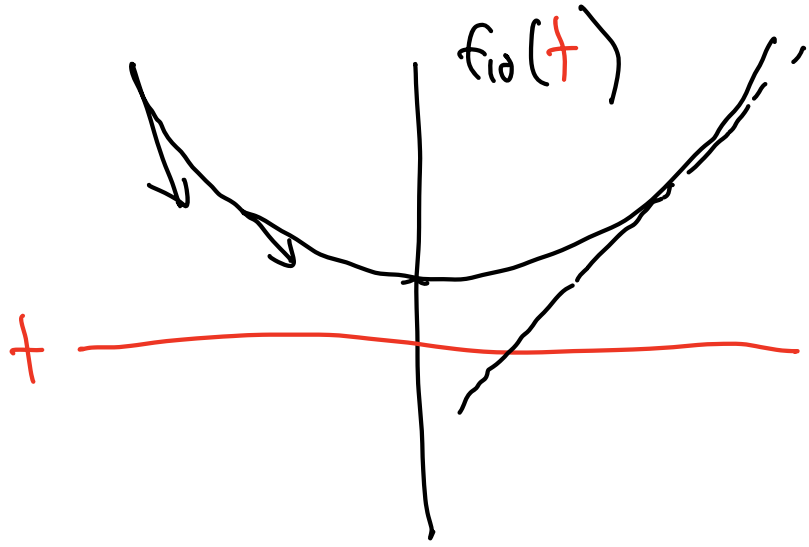
Fact 1:

$$f_{\bar{x},v}(u) \geq f_{\bar{x},v}(t) + f'_{\bar{x},v}(t)(u-t)$$

"above the tangent line"

Fact 2:

$$f''_{\bar{x},v}(t) \geq 0$$



So, what is $f_{10}'(t)$?

$$\frac{\partial}{\partial t} f(\underbrace{\bar{x} + tv}_{x}) = \frac{\partial f}{\partial x_1}(x)\frac{\partial x_1}{\partial t} + \dots$$

$$= \frac{\partial f}{\partial x_1}(x) v_1 + \dots + \frac{\partial f}{\partial x_d}(x) v_d = \nabla f(x)^T v$$

Similarly, $f_{10}''(t) = v^T \underbrace{\nabla^2 f(x)}_{d \times d \text{ matrix of } \frac{\partial^2 f}{\partial x_i \partial x_j}} v$

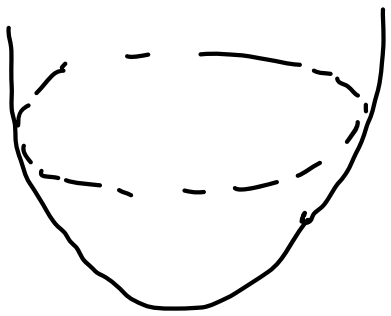
In $\mathbb{R}^d$: f is convex

Fact 1: $f(y) \geq f(x) + \underbrace{\nabla f(x)^T (y - x)}_{\text{the "tangent line"}}$

Fact 2: $v^T \nabla^2 f(x) v \geq 0 \quad \forall v \in \mathbb{R}^d$

Punchline: in $\mathbb{R}^d$, f is convex

$\iff$ $\nabla^2 f(x)$ is PSD



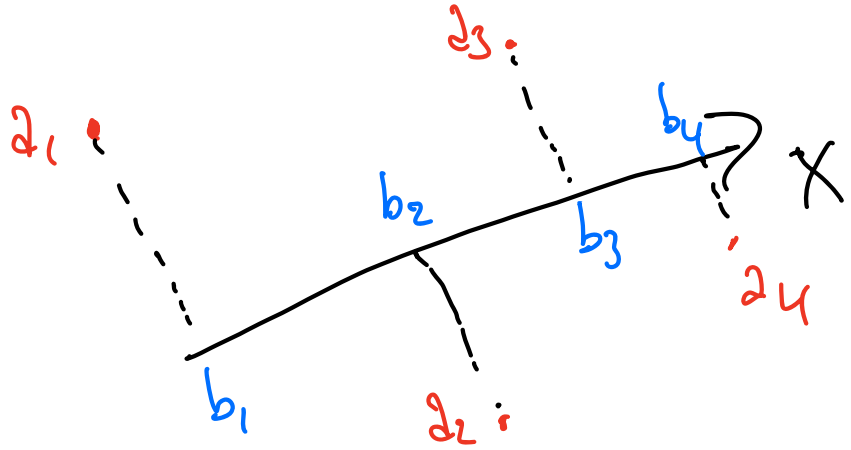
Linear regression (Part VI, Section 5.2)

Input: $\{(a_1, b_1) \dots (a_n, b_n)\}$

$a_i \in \mathbb{R}^d$ "feature vector"

$b_i \in \mathbb{R}$ "response"

Output: X that "explains" relationship



Goal: $\min_{x \in \mathbb{R}^d} \sum_{i \in [n]} \left(a_i^T x - b_i \right)^2$

model of relationship is linear

(HW: piecewise const.)

Another way to write: $A = \begin{pmatrix} a_1^T \\ a_2^T \\ \vdots \\ a_n^T \end{pmatrix}$

$$f(x) = \| Ax - b \|_2^2$$

$$= (Ax - b)^T (Ax - b)$$

$$= x^T \underbrace{A^T A}_{M} x - 2 b^T A x + \text{Const.}$$

Recall: M is PSD.

Fact: if $f(x) = x^T M x + V^T x$

then $\nabla^2 f(x) = 2M$

Regression is convex ☺

For "nice" M we can solve fast using

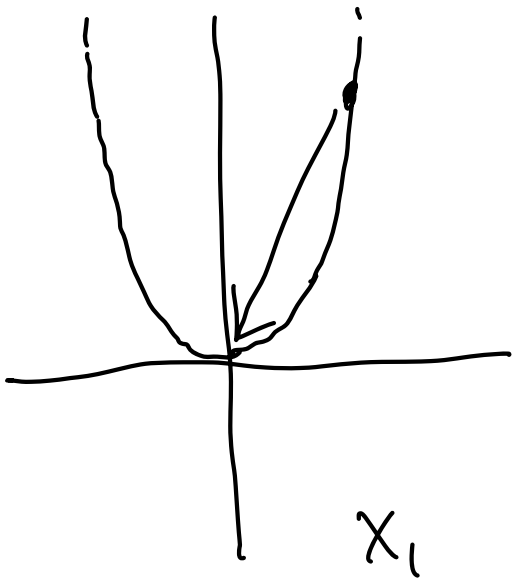
Gradient descent

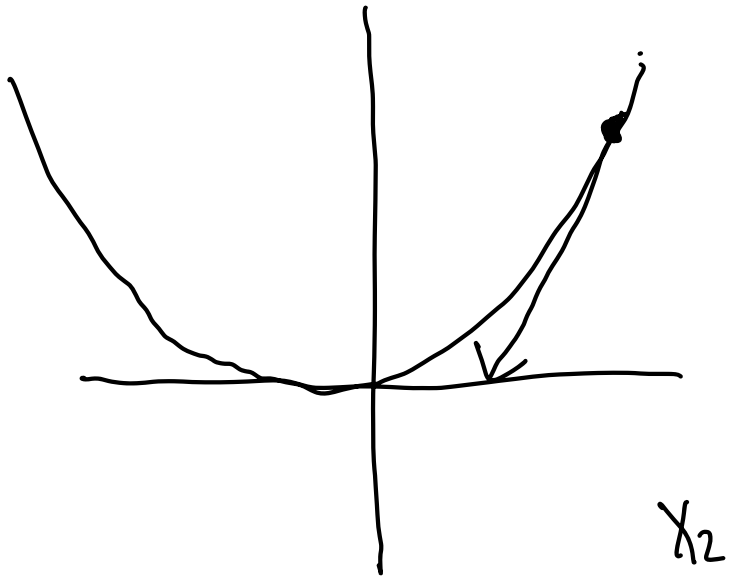
Sketch: Suppose $M = \begin{pmatrix} \lambda_1 & 0 \\ 0 & \lambda_2 \end{pmatrix}$ (diagonal)

$$f(x) = \left(\lambda_1 x_1^2 + V_1 x_1\right) + \left(\lambda_2 x_2^2 + V_2 x_2\right)$$

only get to pick one step size...

$\eta = \frac{1}{2\lambda_1}$ perfect for x_1





perfect

undershot by factor $\frac{\lambda_2}{\lambda_1}$

Takes $\approx \frac{\lambda_1}{\lambda_2}$ iterations to converge.

General case: $M = U \begin{pmatrix} \lambda_1 & & \\ & \lambda_2 & \\ & & \ddots \\ & & & \lambda_0 \end{pmatrix} U^T$

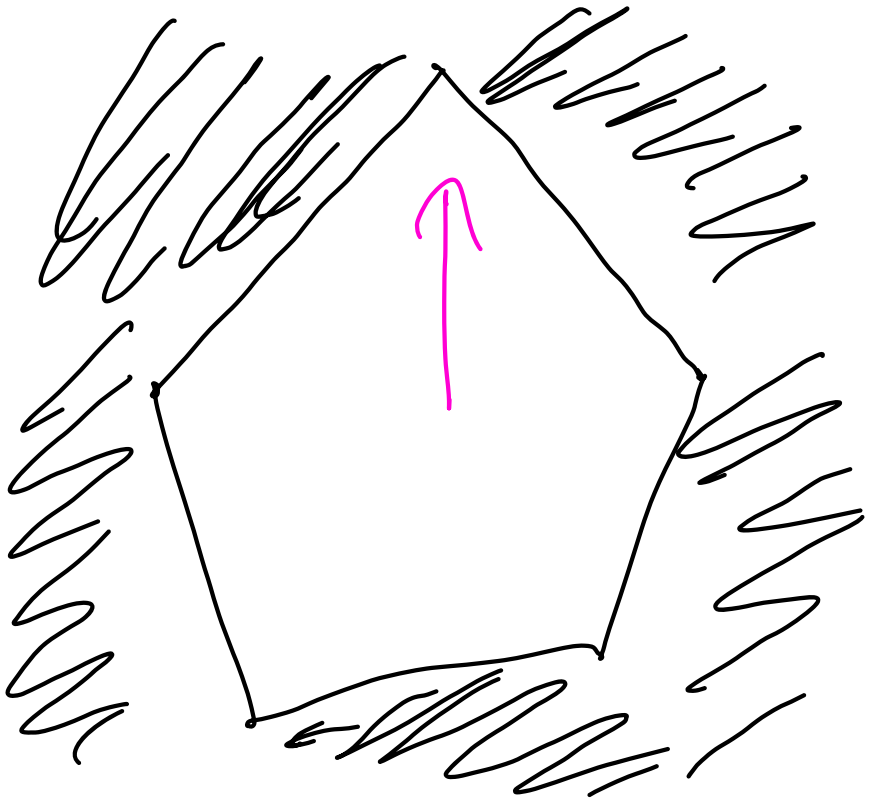
- GD converges in $\approx \frac{\lambda_1}{\lambda_0}$ iteration

"how spherical is M"?

- No problem if $U \neq I$

"input world = output world"

Algos for LP (Part VI, Section 5.3)



It's convex, but...

- constrained
- non-smooth

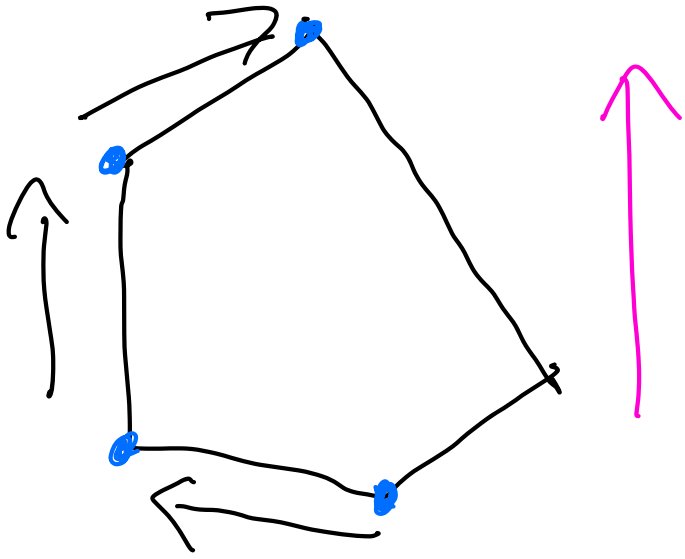
how to solve?

Strongly polynomial: Open "Smale's 9th problem"

Weakly polynomial: ✓ Great in practice!

- 10 most influential of 20th Century
- NYT headline: "Breakthrough in Problem Solving"

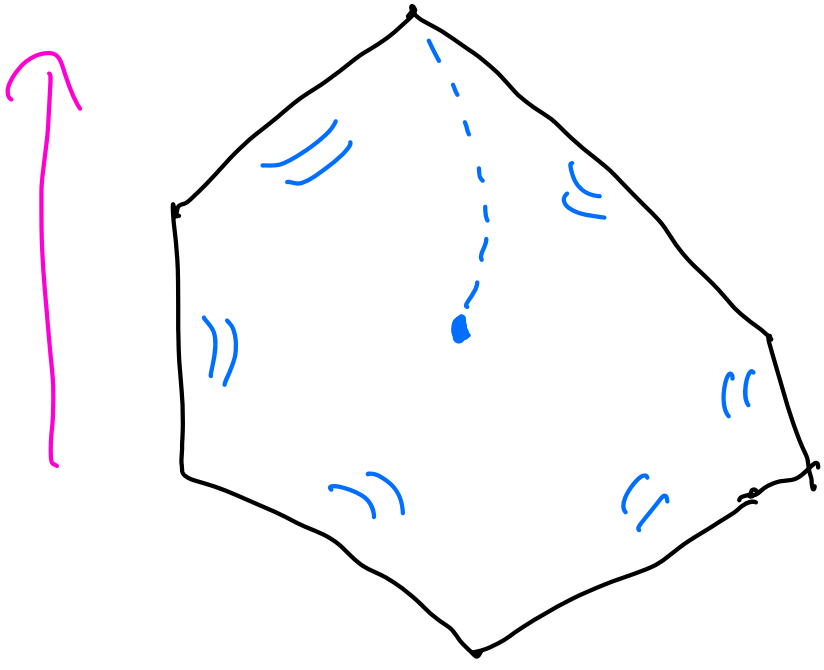
Idea 1: Simplex "walk along the boundary"



- decent in practice
- Worst-case exp time
- "most instances" poly time

Idea 2: IPM "decrease the force fields"

$$x_{t} = \arg\max \; c^{T}x + t\underbrace{b(x)}_{\text{"barrier"}}$$



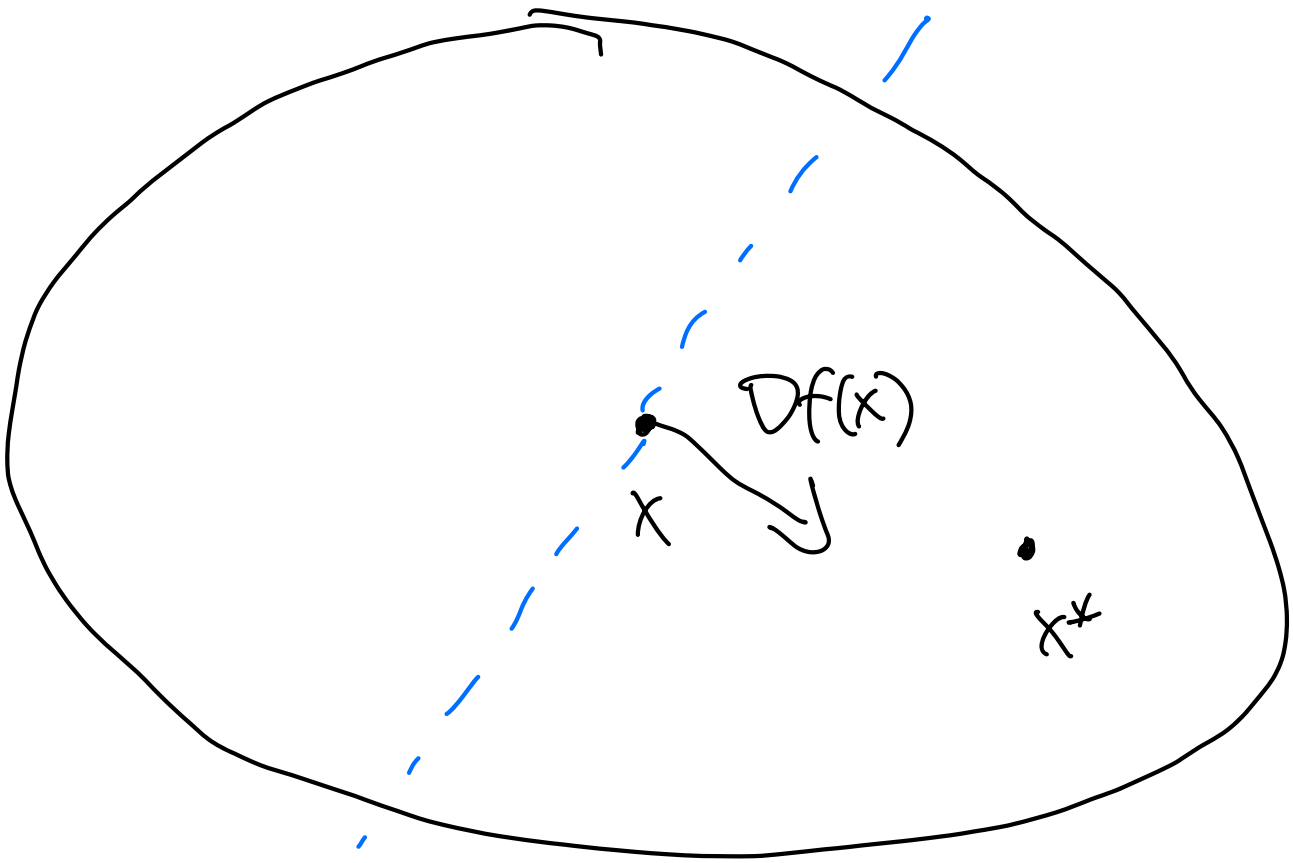
- poly thm
- SOTA in theory
- SOTA (most(-1)) in practice

Idea 3: CPM "binary search in $\mathbb{R}^d$"

- Works for all convex functions
- My favorite also ☺

In 1-d: $f'(x)$ lets you binary search

In $\mathbb{R}^d$: use $\nabla f(x)$



$\nexists$ plane through center-of-gravity x, $\geq 30\%$ on both sides. Decrease volume!